

GAME THEORY-III

(M.Sc. Sem-III)

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GRAPHIC SOLUTION OF $2 \times N$ AND $M \times 2$ GAMES

The procedure described in the last section will generally be applicable for any game with 2×2 payoff matrix unless it possesses a saddle point. Moreover, the procedure can be extended to any square payoff matrix of any order. But it will not work for the game whose payoff matrix happens to be a rectangular one, say $m \times n$. In such cases a very simple graphical method is available if either m or n is two. The graphic short-cut enables us to reduce the original $2 \times n$ or $m \times 2$ game to a much simpler 2×2 game. Consider the following $2 \times n$ game :

$$\begin{array}{c} \text{Player A} \end{array} \left(\begin{array}{cccc} & \text{Player B} & & \\ & B_1 & B_2 & \dots & B_n \\ A_1 & a_{11} & a_{12} & \dots & a_{1n} \\ A_2 & a_{21} & a_{22} & \dots & a_{2n} \end{array} \right)$$

It is assumed that the game does not have a saddle point. Let the optimum mixed strategy for A be given

by $S_A = \begin{bmatrix} p_1 & p_2 \end{bmatrix}$ where $p_1 + p_2 = 1$. The average (expected) payoff for A when he plays S_A against B 's

pure moves B_1, B_2, \dots, B_n is given by

B 's pure move	A 's expected payoff $E(p)$
B_1	$E_1(p_1) = a_{11}p_1 + a_{21}p_2 = a_{11}p_1 + a_{21}(1 - p_1)$
B_2	$E_2(p_1) = a_{12}p_1 + a_{22}p_2 = a_{12}p_1 + a_{22}(1 - p_1)$
\vdots	\vdots
B_n	$E_n(p_1) = a_{1n}p_1 + a_{2n}p_2 = a_{1n}p_1 + a_{2n}(1 - p_1)$

According to the maximum criterion for mixed strategy games, player A should select the values of p_1 and p_2 so as to maximize his minimum expected payoffs. This may be done by plotting the expected payoff lines :

$$E_j(p_1) = (a_{1j} - a_{2j})p_1 + a_{2j} \quad (j = 1, 2, \dots, n).$$

The highest point on the lower envelope of these lines will give maximum of the minimum (i.e., maximin) expected payoffs to player A as also the maximum value of p_1 .

The two lines* passing through the maximin point identify the two critical moves of B which combined with two of A , yield the 2×2 matrix that can be used to determine the optimum strategies of the two players, for the original game, using the results of the previous section.

The $(m \times 2)$ games are also treated in the same way where the upper envelope of the straight lines corresponding to B 's expected payoffs will give the maximum expected payoff to player B and the lowest point on this then gives the minimum expected payoff (minimax value) and the optimum value of q_1 .

SAMPLE PROBLEMS

1. Solve the following 2×2 game graphically :

		Player B			
		B_1	B_2	B_3	B_4
Player A	A_1	2	1	0	-2
	A_2	1	0	3	2

Solution. Clearly, the problem does not possess a saddle point. Let the player A play the mixed strategy

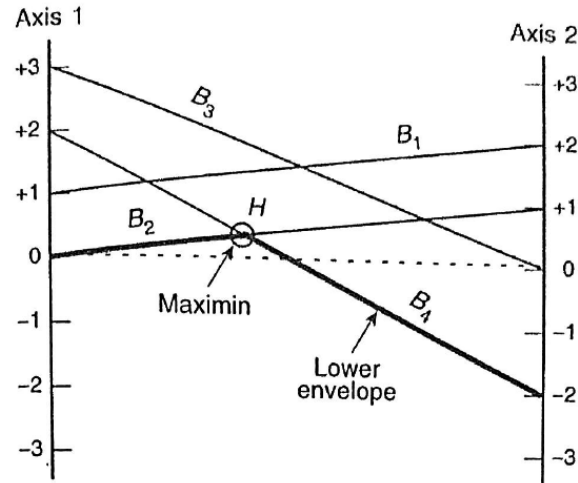
$$S_A = \begin{bmatrix} A_1 & A_2 \\ p_1 & p_2 \end{bmatrix} \text{ where } p_2 = 1 - p_1, \text{ against } B. \text{ Then } A\text{'s expected payoffs against } B\text{'s pure moves are given}$$

by

B 's pure move	A 's expected payoff $E(p_1)$
B_1	$E_1(p_1) = p_1 + 1$
B_2	$E_2(p_1) = p_1$
B_3	$E_3(p_1) = -3p_1 + 3$
B_4	$E_4(p_1) = -4p_1 + 2$

These expected payoff equations are then plotted as functions of p_1 as shown in figure. Which shows the payoffs of each column represented as points on two vertical axis 1 and 2, unit distance apart. Thus line B_1 joins the first payoff element 2 in the first column represented by +2 on axis 2 and the second payoff element 1 in the first column represented by +1 on axis 1. Similarly, lines B_2 , B_3 and B_4 join the corresponding representation of payoff elements in the second, third and fourth columns. Since the player A wishes to maximize his minimum expected payoff we consider the highest point of intersection H on the lower envelope of the A 's expected payoff equations. This point H represents the maximum expected value of the game for A . The lines B_2 and B_4 , passing through H , define the two relevant moves B_2 and B_4 that alone B needs to play. The solution to the original 2×4 game, therefore, boils down that of the simpler game with the 2×2 payoff matrix :

	B_2	B_4
A_1	1	-2
A_2	0	2



The maximin value

Now if $S_A = \begin{bmatrix} A_1 & A_2 \\ p_1 & p_2 \end{bmatrix}$ and $S_B = \begin{bmatrix} B_2 & B_4 \\ q_2 & q_4 \end{bmatrix}$

be the optimum strategies for A and B , then we have

$$p_1 = \frac{2-0}{1+2-(-2)} = 2/5, \quad p_2 = 1 - p_1 = 3/5,$$

$$q_2 = \frac{2-(-2)}{1+2-(-2)} = 4/5, \quad q_4 = 1 - q_2 = 1/5.$$

Hence, the solution to the game is

(i) the optimum strategy for A is $S_A = \begin{bmatrix} A_1 & A_2 \\ 2/5 & 3/5 \end{bmatrix}$

(ii) the optimum strategy for B is : $S_B = \begin{bmatrix} B_1 & B_2 & B_3 & B_4 \\ 0 & 4/5 & 0 & 1/5 \end{bmatrix}$

and (iii) the expected value of the game is $v = \frac{2 \times 1 - 0 \times (-2)}{1 + 2 - (0 - 2)} = \frac{2}{5}.$

2. Obtain the optimal strategies for both-persons and the value of the game for zero-sum two person game whose payoff matrix is as follows :

$$\begin{bmatrix} 1 & -3 \\ 3 & 5 \\ -1 & 6 \\ 4 & 1 \\ 2 & 2 \\ -5 & 0 \end{bmatrix}$$

Sol. Clearly, the given problem does not possess any saddle point. So, let the player B play the mixed strategy

$$S_B = \begin{bmatrix} B_1 & B_2 \\ q_1 & q_2 \end{bmatrix} \text{ with } q_2 = 1 - q_1 \text{ against player } A. \text{ Then } B\text{'s expected payoffs against } A\text{'s pure moves are}$$

given by

A 's pure move

B 's expected payoff $E(q_1)$

A_1

$$E_1(q_1) = 4q_1 - 3$$

A_2

$$E_2(q_1) = -2q_1 + 5$$

A_3

$$E_3(q_1) = -7q_1 + 6$$

A_4

$$E_4(q_1) = 3q_1 + 1$$

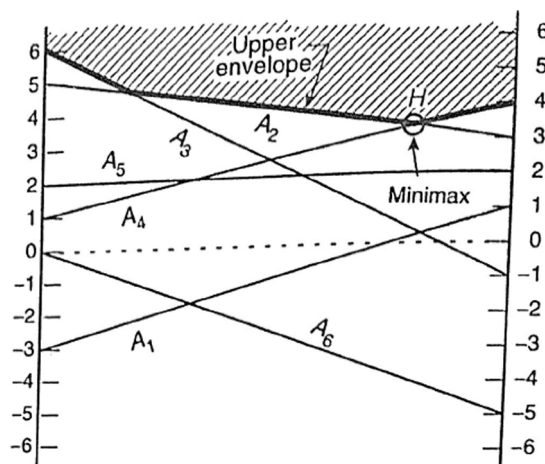
A_5

$$E_5(q_1) = 2$$

A_6

$$E_6(q_1) = -5q_1$$

The expected payoff equations are then plotted as functions of q_1 as shown in figure.



The minimax value

Since, the player B wishes to minimize his maximum expected payoff, we consider the loose point of intersection H on the upper envelope of B 's expected payoff equations. This point B represents the minimax expected value of the game for player B . The lines A_2 and A_4 passing through H , define the two relevant moves A_2 and A_4 that alone the player A needs to play. The solution to the original 6×2 game, therefore,

reduces to that of the simpler game with 2×2 payoff matrix :

$$\begin{array}{c} \text{Player B} \\ \text{Player A} \end{array} \begin{bmatrix} 3 & 5 \\ 4 & 1 \end{bmatrix}$$

If we now, let

$$S_A = \begin{bmatrix} A_1 & A_2 \\ p_1 & p_2 \end{bmatrix}, p_1 + p_2 = 1; S_B = \begin{bmatrix} B_1 & B_2 \\ q_1 & q_2 \end{bmatrix}, q_1 + q_2 = 1$$

then using the usual method of solution for 2×2 games, the optimum strategies can easily be obtained as

$$S_A = \begin{bmatrix} A_1 & A_2 & A_3 & A_4 & A_5 & A_6 \\ 0 & 3/5 & 0 & 2/5 & 0 & 0 \end{bmatrix}, S_B = \begin{bmatrix} B_1 & B_2 \\ 4/5 & 1/5 \end{bmatrix}$$

and the value of the game as $v = 17/5$.